

Exercises for FFL, Fall 2017, sheet 8

Return Thursday Nov 23, in class.

Exercise 1. Let $S = \{c, f\}$ be a signature consisting of a constant symbol c and a binary function symbol f . Let $\mathcal{A} = (A, c^{\mathcal{A}}, f^{\mathcal{A}})$ be the S -structure over carrier $A = \{1, 2\}$ with $c^{\mathcal{A}} = 1$ and $f^{\mathcal{A}}$ given by $f^{\mathcal{A}}(1, 1) = 1, f^{\mathcal{A}}(1, 2) = 2, f^{\mathcal{A}}(2, 1) = f^{\mathcal{A}}(2, 2) = 1$. Your task: represent \mathcal{A} as a set, coding lists by iterated nested pairs (right-associative) and representing these (nested) pairs as sets (so ultimately there will be only “{ }” symbols, no more “()” symbols). Using a text processor’s copy-paste capabilities saves a lot of time in this task, prevents copy errors, and makes it possible for the TAs to read your result, so please use a text processor for this task.

Solution. Let us work from the inside to the outside. The graph set of $f^{\mathcal{A}}$, written using pairs, is $\{((1,1),1), ((1,2),2), ((2,1),1), ((2,2),2)\}$. Turning the argument pairs into sets turns this to $\{(\{1,\{1\}\},1), (\{1,\{1,2\}\},2), (\{2,\{1,2\}\},1), (\{2,\{2\}\},2)\}$ and the outer pair brackets into sets to $f^{\mathcal{A}} = \{\{\{1,\{1\}\},\{\{1,\{1\}\},1\}\}, \{\{1,\{1,2\}\},\{\{1,\{1,2\}\},2\}\}, \{\{2,\{1,2\}\},\{\{2,\{1,2\}\},1\}\}, \{\{2,\{2\}\},\{\{2,\{2\}\},2\}\}\}$.

$c^{\mathcal{A}}$ is 1.

The pair $(c^{\mathcal{A}}, f^{\mathcal{A}})$ is $\{c^{\mathcal{A}}, \{c^{\mathcal{A}}, f^{\mathcal{A}}\}\} = \{1, \{1, \{\{\{1,\{1\}\},\{\{1,\{1\}\},1\}\}, \{\{1,\{1,2\}\},\{\{1,\{1,2\}\},2\}\}, \{\{2,\{1,2\}\},\{\{2,\{1,2\}\},1\}\}, \{\{2,\{2\}\},\{\{2,\{2\}\},2\}\}\}\}$.

The triple $(A, c^{\mathcal{A}}, f^{\mathcal{A}})$ then becomes $(A, (c^{\mathcal{A}}, f^{\mathcal{A}})) = \{A, \{A, (c^{\mathcal{A}}, f^{\mathcal{A}})\}\} = \{\{1,2\}, \{\{1,2\}, \{1, \{1, \{\{\{1,\{1\}\},\{\{1,\{1\}\},1\}\}, \{\{1,\{1,2\}\},\{\{1,\{1,2\}\},2\}\}, \{\{2,\{1,2\}\},\{\{2,\{1,2\}\},1\}\}, \{\{2,\{2\}\},\{\{2,\{2\}\},2\}\}\}\}\}$.

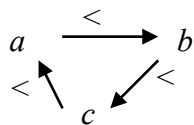
Exercise 2.

Let $S = \{<\}$, where $<$ is a binary relation symbol. Characterize in words the class of all S -structures \mathcal{A} which are models of

$$\varphi = \forall x_1 \forall x_2 \forall x_3 (((\neg x_1 = x_2 \wedge \neg x_2 = x_3) \wedge \neg x_1 = x_3) \wedge < x_1 x_2) \wedge < x_2 x_3) \rightarrow \neg < x_3 x_1).$$

Furthermore, give two concrete S -structures, one of which is a model of φ and the other isn't. Present your structures by specifying the carrier and the extension of $<^{\mathcal{A}}$.

Solution. The models of φ are exactly those $\{<\}$ -structures that contain no $<$ -cycle of length 3. The simplest S -structure that is a model of φ is given by a singleton set A and empty $<$, that is, $A = \{a\}$ and $<^{\mathcal{A}} = \emptyset$ (graph-like representation: a single point). The simplest S -structure that is not a model of φ is an isolated 3-cycle of $<$, that is, $A = \{a, b, c\}$ and $<^{\mathcal{A}} = \{(a,b), (b,c), (c,a)\}$:



Exercise 3. Let $S = \{\text{father-of}, \text{Tom}, \text{Anne}\}$, where father-of is a unary function symbol and Tom, Anne are constant symbols. Give a concrete S -structure which is *not* a model of the S -

expression "father-of Anne = Tom". And give a concrete S -structure which is a model of the S -expression "father-of Tom = Anne".

Solution (out of a zillion possibilities).

(a) $\mathcal{A} = (A, \text{Tom}^{\mathcal{A}}, \text{Anne}^{\mathcal{A}}, \text{father-of}^{\mathcal{A}}) = (\{1, 2\}, 1, 2, \{(1, 1), (2, 2)\})$

(b) $\mathcal{A} = (A, \text{Tom}^{\mathcal{A}}, \text{Anne}^{\mathcal{A}}, \text{father-of}^{\mathcal{A}}) = (\{1\}, 1, 1, \{(1, 1)\})$

Exercise 4. You know that a binary relation is called an *equivalence relation* if it is (i) reflexive, (ii) symmetric, and (iii) transitive. Let $S = \{\equiv\}$, where \equiv is a binary relation symbol. Give an S -expression ϕ such that in any S -structure $\mathcal{A} = (A, \equiv^{\mathcal{A}})$, which is a model of ϕ , $\equiv^{\mathcal{A}}$ is an equivalence relation on A .

Solution. The following S -expression characterizes $\equiv^{\mathcal{A}}$ to be an equivalence relation:

$$\forall x (\equiv x x) \wedge \forall x \forall y (\equiv x y \rightarrow \equiv y x) \wedge \forall x \forall y \forall z ((\equiv x y \wedge \equiv y z) \rightarrow \equiv x z)$$